2:23 Part 2 Kevin Brown

Let schedule $S_1 \longrightarrow S_2 : f_1(x); f_2(z); f_3(x); f_4(z); f_5(y); f_6(y); w_1(x); w_2(z); w_3(y); w_2(y)$

From the order, $T_3$ interferes in execution of $T_2$, which makes schedule non-serializable

- $T_2$ reads $Y(f_2(x))$, which is read and modified by $T_3(w_3(x))$
- $T_3$ reads $Y(f_5(y))$, which is modified before $T_2$ modifies $Y(w_2(y))$